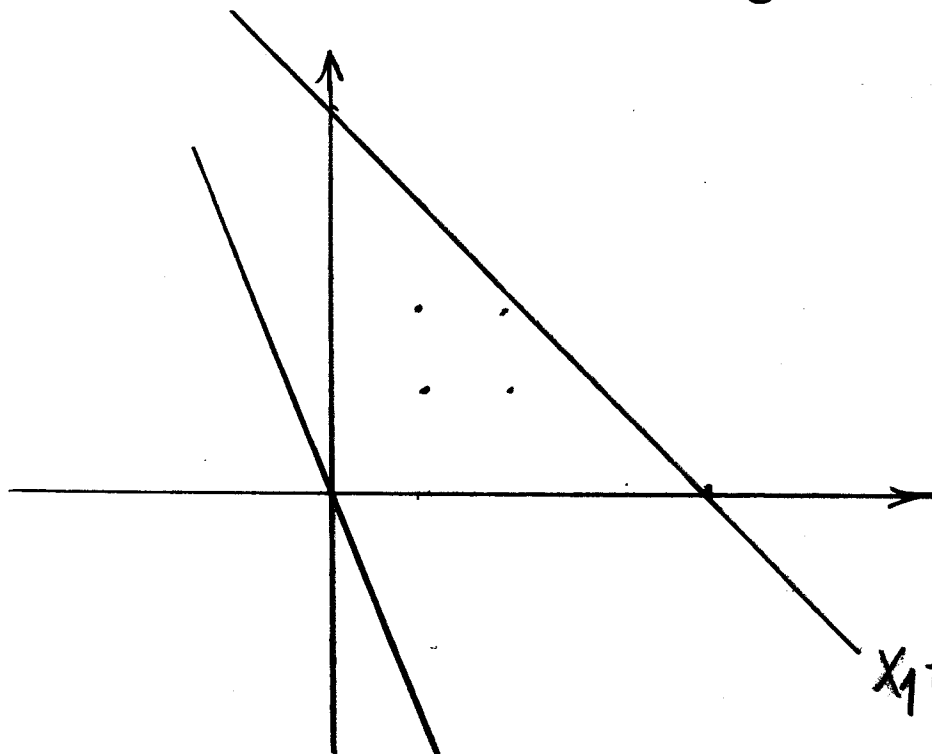


Integer Programming

$$\begin{aligned} \max \quad & 3x_1 + 2x_2 \\ & x_1 + x_2 \leq 6 \\ & x_1 \geq 0, \quad x_2 \geq 0 \end{aligned}$$

LP relaxation for Integer Programming

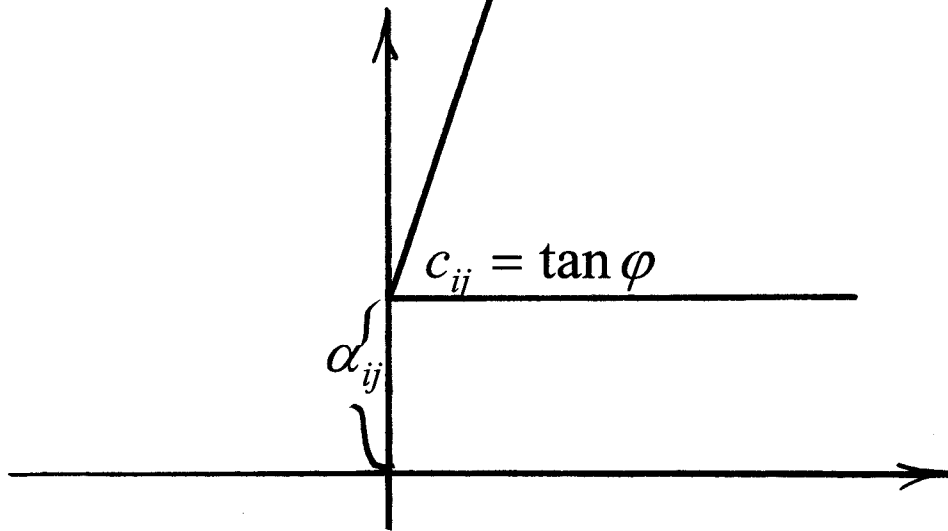


$$z = 3x_1 + 2x_2$$

$$\begin{aligned} 3x_1 + 2x_2 &\rightarrow \max \\ x_1 + x_2 &\leq 6 \\ 4x_1 + x_2 &\leq 8 \\ x_1 \geq 0, \quad x_2 &\geq 0 \end{aligned}$$

Fixed Charge Transportation

$$\sum \sum c_{ij}(x_{ij}) \rightarrow \min$$



$$c_{ij}(x_{ij}) = \begin{cases} c_{ij}x_{ij} + \alpha_{ij}, & x_{ij} > 0 \\ 0, & x_{ij} = 0 \end{cases}$$

$$m_{ij} = \min\{a_i, b_j\}$$

$$\sum \sum c_{ij} x_{ij} + \sum \sum \alpha_{ij} y_{ij} \rightarrow \min$$

$$\sum_j x_{ij} = a_i, \quad \sum_i x_{ij} = b_j$$

$$0 \leq x_{ij} \leq m_{ij} y_{ij}$$

$$y_{ij} \in \{0, 1\}$$

	<i>Shirts</i>	<i>Shorts</i>	<i>Pants</i>	
<i>Labor</i>	3	2	6	150
<i>Cloth</i>	4	3	4	100

	<i>Shirts</i>	<i>Shorts</i>	<i>Pants</i>
<i>Sales Pr.</i>	\$12	\$8	\$15
<i>Var. Cost</i>	\$6	\$4	\$8

Machinery I - \$200, II - \$150, III - \$100

Shirts - x_1 , *Shorts* - x_2 , *Pants* - x_3

Mach I - y_1 , *Mach II* - y_2 , *Mach III* - y_3

$$z = 6x_1 + 4x_2 + 7x_3 - 200y_1 - 150y_2 - 100y_3$$

$$3x_1 + 2x_2 + 6x_3 \leq 150$$

$$4x_1 + 3x_2 + 4x_3 \leq 100$$

$$x_1 \geq 0, x_2 \geq 0, x_3 \geq 0$$

$$y_i \in \{0, 1\}$$

$$x_1 \leq \mu_1 y_1$$

$$x_2 \leq \mu_2 y_2$$

$$x_3 \leq \mu_3 y_3$$

Either or Constraints

Dorian Auto is considering manufacturing three types of autos: compact, midsize & large

6000 tons of steel

60000 hrs. of labor is available

$$\begin{aligned} \text{Constraints } x_1 &\leq 0 \quad \text{or} \quad x_1 \geq 1000 \\ x_2 &\leq 0 \quad \text{or} \quad x_2 \geq 1000 \\ x_3 &\leq 0 \quad \text{or} \quad x_3 \geq 1000 \end{aligned}$$

	<i>Comp.</i>	<i>Midsize</i>	<i>Large</i>
<i>Steel</i>	1.5T	3T	5T
<i>Labor</i>	30hr	25hr	40hr
<i>Profit yielded</i>	\$2000	\$3000	\$4000

$$\begin{array}{ll} x_1 - \text{comp. cars} & x_1 \leq 2000 \\ x_2 - \text{midsize} & x_2 \leq 2000 \\ x_3 - \text{large} & x_3 \leq 1200 \end{array}$$

$$\begin{array}{ll} & x_1 \leq 2000y_1 \\ 1000 - x_1 \leq 2000(1 - y_1) & y_1 \in \{0, 1\} \end{array}$$

$$\begin{array}{ll} & x_2 \leq 2000y_2 \\ 1000 - x_2 \leq 2000(1 - y_2) & y_2 \in \{0, 1\} \end{array}$$

$$\begin{array}{ll} & x_3 \leq 1200y_3 \\ 1000 - x_3 \leq 1200(1 - y_3) \end{array}$$

$$1.5x_1 + 3x_2 + 5x_3 \leq 6000$$

$$30x_1 + 25x_2 + 40x_3 \leq 60000$$

$$x_1 \geq 0, x_2 \geq 0, x_3 \geq 0 \quad x_i - \text{integer}$$

$$y_i \in \{0, 1\}$$

$$z = 2x_1 + 3x_2 + 4x_3$$

$$x_1 \leq 2000y_1$$

$$1000 - x_1 \leq 2000(1 - y_1)$$

$$x_2 \leq 2000y_2$$

$$1000 - x_2 \leq 2000(1 - y_2)$$

$$x_3 \leq 1200y_3$$

$$1000 - x_3 \leq 1200(1 - y_3)$$

$$1.5x_1 + 3x_2 + 5x_3 \leq 6000$$

$$30x_1 + 25x_2 + 40x_3 \leq 60000$$

$$x_i \geq 0, \quad y_i \in \{0, 1\}$$

If - Then Constraints

If $f(x_1, \dots, x_n) > 0$ then $g(x_1, \dots, x_n) \leq 0$ must be satisfied

$$\begin{aligned} -g &\leq \mu y \\ f &\leq \mu(1-y) \end{aligned}$$

If $f > 0$ then only $y = 0$ is possible, so

$$-g \leq 0 \text{ or } g \geq 0$$

Set Covering

		<i>Math1</i>	<i>Math2</i>	<i>OR1</i>	<i>OR2</i>	<i>CS1</i>	<i>CS2</i>
x_1	<i>Calculus</i>	1	1				
x_2	<i>OR</i>	1	1	1	1	1	
x_3	<i>Data Struct</i>	1				1	1
x_4	<i>Business Stat</i>	1	1	1			
x_5	<i>Comp. Simul</i>				1	1	
x_6	<i>Introd. to CP</i>						1
x_7	<i>Forecasting</i>	1		1	1		

Cal \Rightarrow *Bus. Stat.*

Math1 - 2

Math2 - 2

Intr. \Rightarrow *Comp. Simul*

OR1 - 1

OR2 - 2

Intr. \Rightarrow *Business Stat.*

CS1 - 1

CS2 - 1

Bus. Stat \Rightarrow *Forecast*

$$\sum x_i \rightarrow \min$$

$$x_1 + x_2 + x_3 + x_4 + x_7 \geq 2$$

$$x_1 \geq x_4$$

$$x_1 + x_2 + x_4 \geq 1$$

$$x_6 \geq x_5$$

$$x_2 + x_4 + x_7 \geq 1$$

$$x_6 \geq x_4$$

$$x_2 + x_5 + x_7 \geq 2$$

$$x_4 \geq x_7$$

$$x_2 + x_3 + x_5 \geq 1$$

$$x_3 + x_6 \geq 1, \quad x_i \in \{0, 1\}$$

	<i>Math1</i>	<i>Math2</i>	<i>OR1</i>	<i>OR2</i>	<i>CS1</i>	<i>CS2</i>	
<i>Calculus</i>	1	1					x_1
<i>OR</i>	1	1	1	1	1		x_2
<i>Data Struct</i>	1				1	1	x_3
<i>Bus. Stat</i>	1	1	1				x_4
<i>Comp. Simul</i>				1	1	1	x_5
<i>Int. CP</i>					1		x_6
<i>Forecasting</i>	1		1	1			x_7
	1	2	2	1	2	1	

$$\sum x_i \rightarrow \min$$

$$x_1 + x_2 + x_3 + x_4 + x_7 \geq 1$$

$$x_1 + x_2 + x_4 \geq 2$$

$$x_2 + x_4 + x_7 \geq 2$$

$$x_2 + x_5 + x_7 \geq 1$$

$$x_2 + x_3 + x_5 + x_6 \geq 2$$

$$x_3 + x_5 \geq 1$$

$$x_i \in \{0, 1\}$$

Set Covering

0	10	20	30	30	20
10	0	25	35	20	10
20	25	0	15	30	20
30	35	15	0	15	25
30	20	30	15	0	14
20	10	20	25	14	0

x_1	1	1				
x_2	1	1			1	
x_3			1	1		
x_4			1	1	1	
x_5				1	1	1
x_6		1			1	1

$$z = x_1 + x_2 + \dots + x_6$$

$$x_3 + x_4 + x_5 \geq 1 \qquad x_1 + x_2 \geq 1$$

$$x_4 + x_5 + x_6 \geq 1 \qquad x_1 + x_2 + x_6 \geq 1$$

$$x_2 + x_5 + x_6 \geq 1 \qquad x_3 + x_4 \geq 1$$

$$x_i \in \{0, 1\}$$

Knapsack Problem

Branch & Bound

$$5x_1 + 8x_2 + 3x_3 + 7x_4 \rightarrow \max$$

$$3x_1 + 5x_2 + 2x_3 + 4x_4 \leq 10$$

$$x_i \in \{0, 1\}$$

$$\frac{7}{4} > \frac{5}{3} > \frac{8}{5} > \frac{3}{2}$$

$$7x_1 + 5x_2 + 8x_3 + 3x_4 \rightarrow \max$$

$$4x_1 + 3x_2 + 5x_3 + 2x_4 \leq 10$$

$$x_i \in \{0, 1\}$$

Relaxation LP

$$7x_1 + 5x_2 + 8x_3 + 3x_4 \rightarrow \max$$

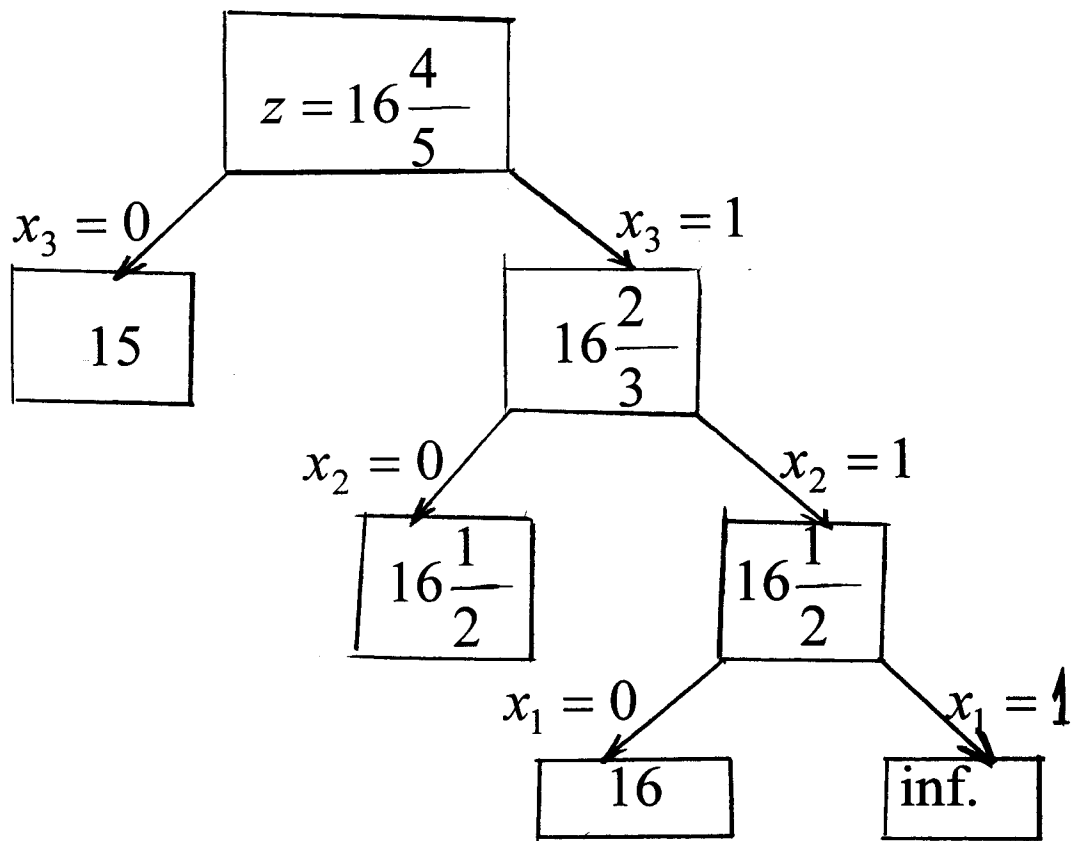
$$4x_1 + 3x_2 + 5x_3 + 2x_4 \leq 10$$

$$0 \leq x_i \leq 1$$

Solution:

$$x_1 = 1, x_2 = 1, x_3 = \frac{3}{5}$$

$$z = 7 + 5 + \frac{3 \cdot 8}{5} = 16\frac{4}{5} \quad \text{upper bound}$$

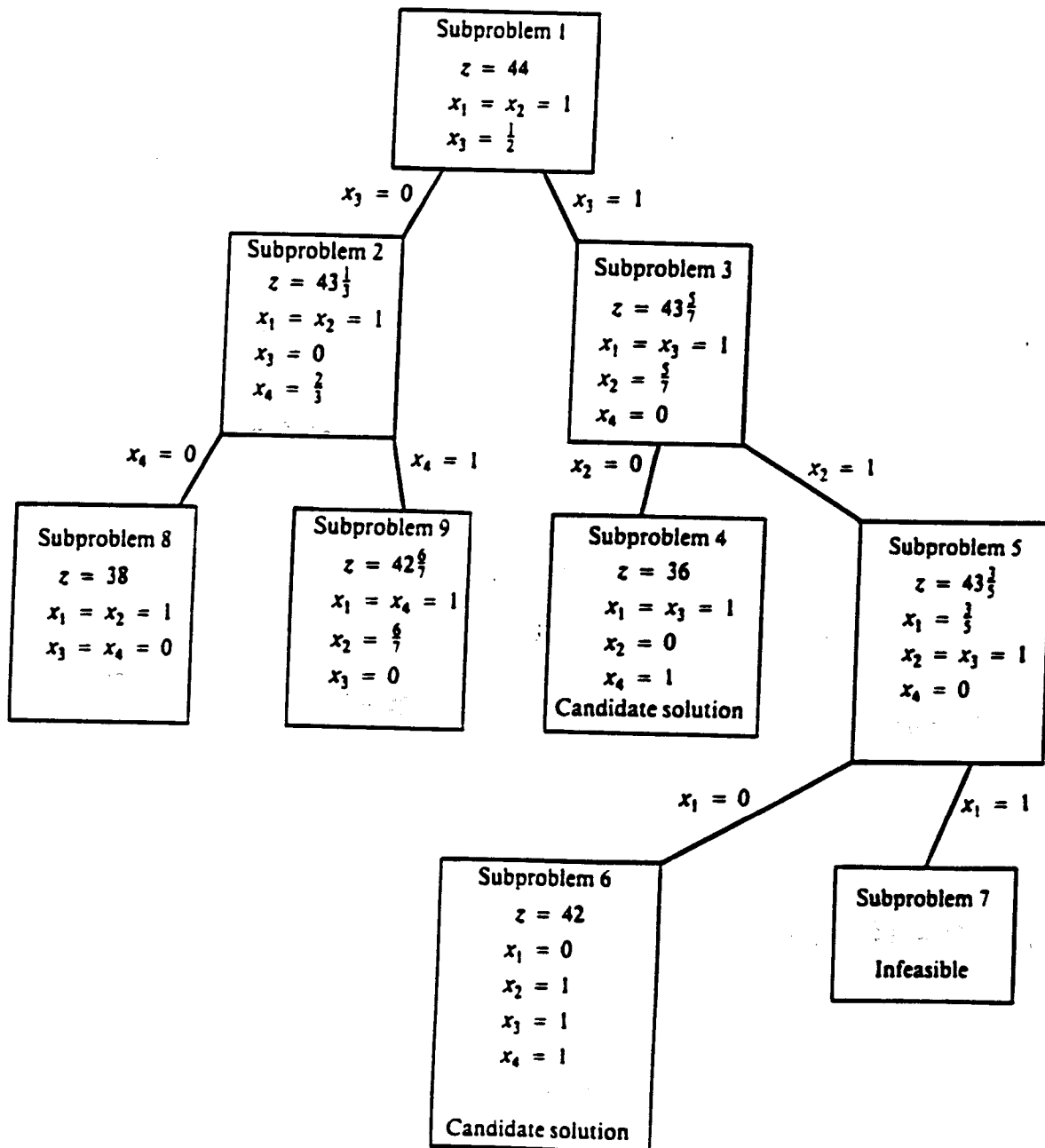


$$z^* = 16, \quad x_1^* = 0, \quad x_2^* = x_3^* = x_4^* = 1$$

$$16x_1 + 22x_2 + 12x_3 + 8x_4 \rightarrow \max$$

$$5x_1 + 7x_2 + 4x_3 + 3x_4 \leq 14$$

$$x_i \in \{0, 1\}$$

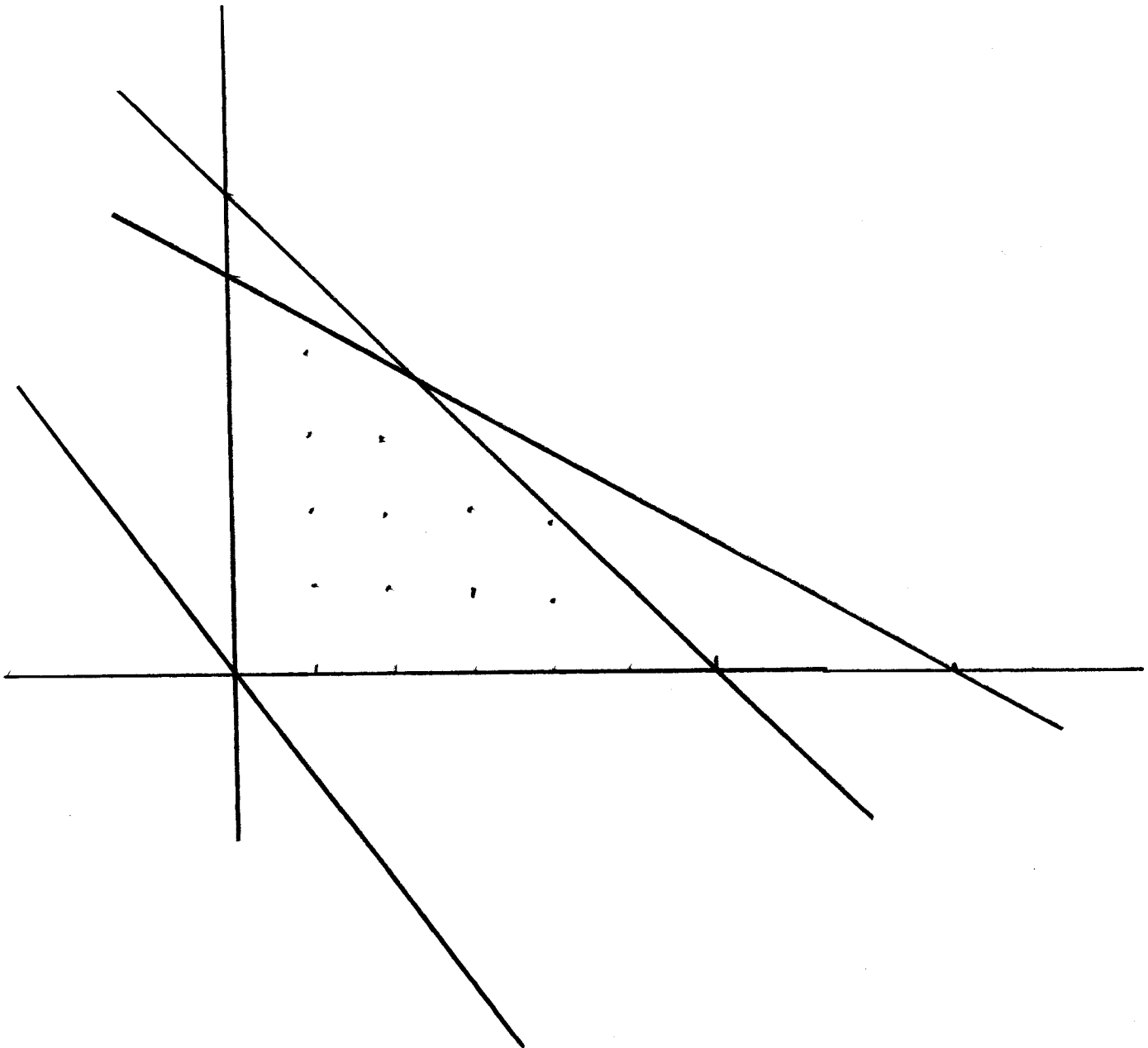


$$z = 8x_1 + 5x_2 \rightarrow \max$$

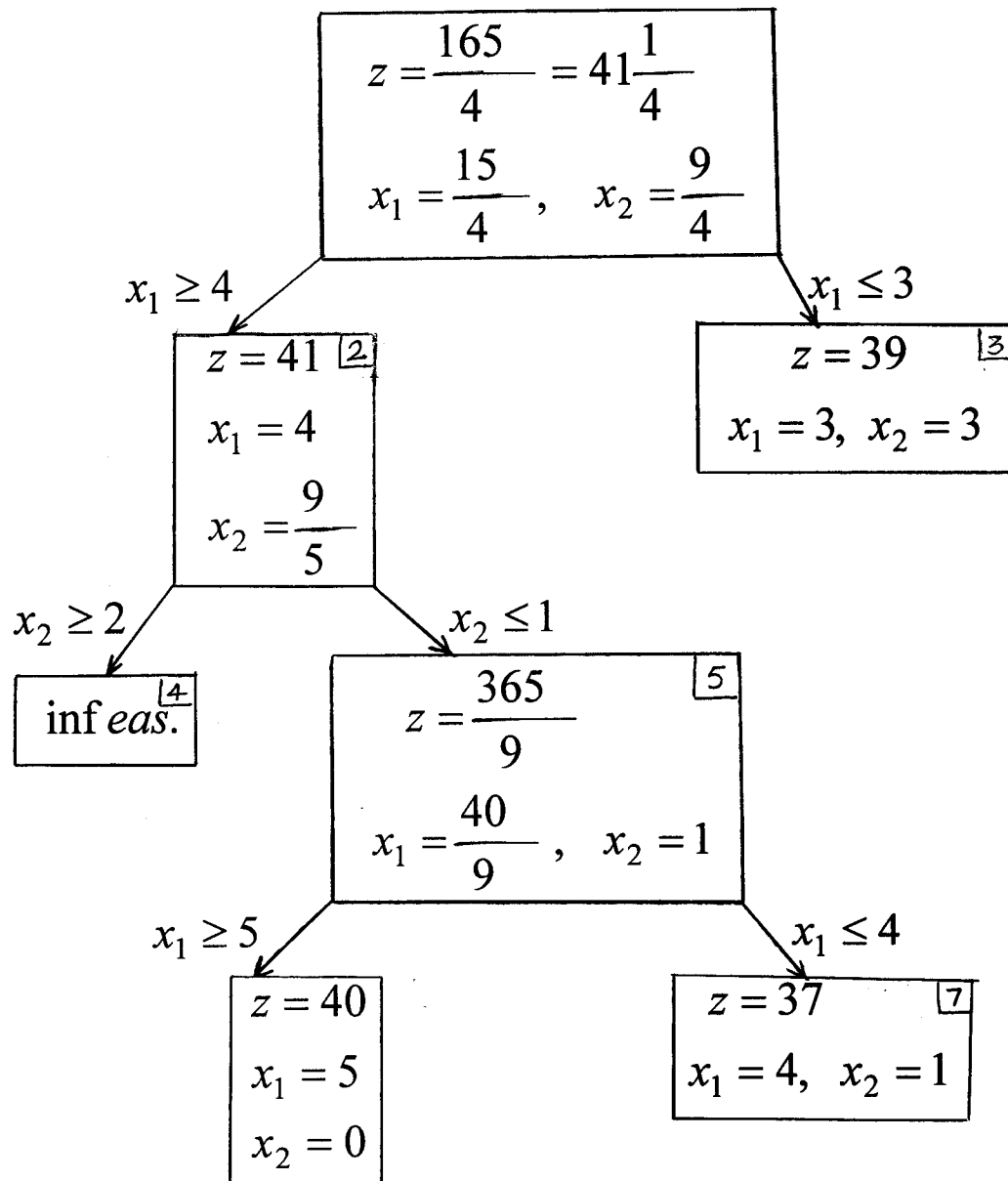
$$x_1 + x_2 \leq 6$$

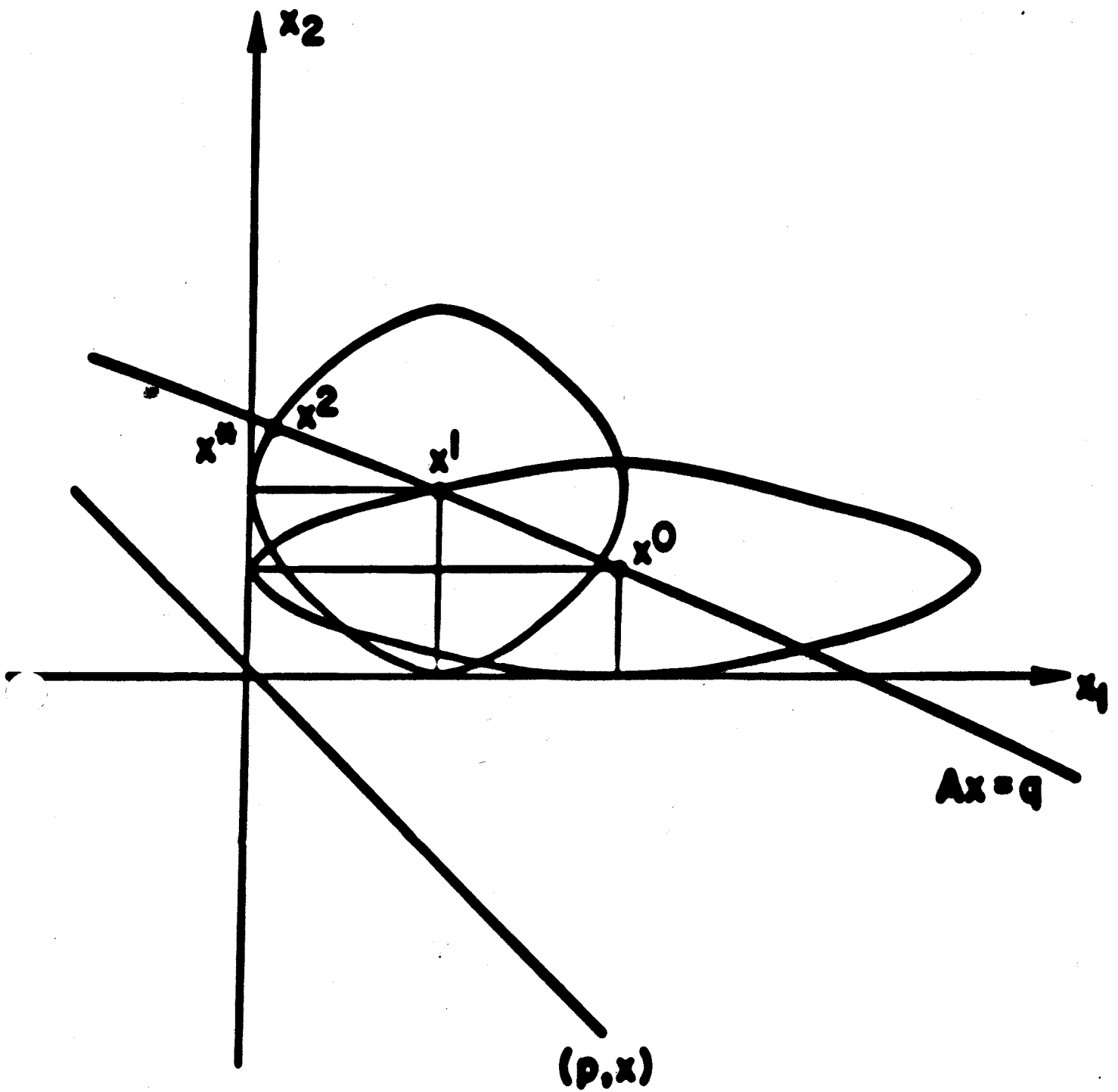
$$9x_1 + 5x_2 \leq 45$$

$$x_1 \geq 0, x_2 \geq 0; x_i - \text{integer}$$



$$z = \frac{165}{4}, \quad x_1 = \frac{15}{4}, \quad x_2 = \frac{9}{4}$$





$$(p, x^{s+1} - x^*) \leq \left(1 - \frac{\alpha}{\sqrt{n-m} + \epsilon_s}\right) (p, x^s - x^*), \quad \epsilon_s \rightarrow 0$$

I.I. Dikin 1967